

# Tournament Matrices

Immanuel Chen

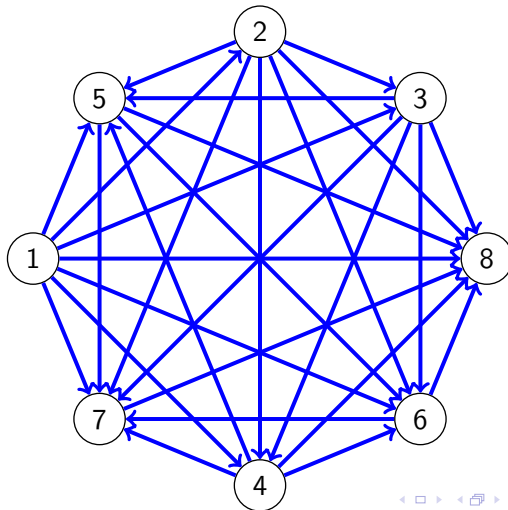
# Tournaments

- ▶ Digraph that represents the outcome of a round-robin tournament

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- ▶ Vertices are teams
- ▶ Edges denotes the victor between two teams
- ▶  $K_n$  with direction

# Example



# Tournament Matrices

- ▶ Adjacency matrix of a tournament

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- ▶ 1's represent wins, 0's represent losses

# Properties

Let  $A$  be a tournament matrix of size  $n \times n$

- ▶  $[A]_{ii} = 0$  for  $1 \leq i \leq n$
- ▶  $[A]_{ij} + [A]_{ji} = 1$  for  $1 \leq i < j \leq n$
- ▶  $A + A^T = J_n - I_n$
- ▶  $\sum_{i=1}^n \sum_{j=1}^n [A]_{ij} = \binom{n}{2}$

## Row and Column Sums

- ▶ Row sum vector  $R = (r_1, r_2, \dots, r_n)$  where  $r_i = \sum_{j=1}^n [A]_{ij}$ 
  - ▶  $r_i$  represents the number of wins team  $i$  has
  - ▶ Also known as the score vector
  - ▶  $R = R(A) = A j_n$



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- ▶ Column sum vector  $S = (s_1, s_2, \dots, s_n)$  where  $s_i = \sum_{j=1}^n [A]_{ij}$ 
  - ▶  $s_i$  represents the number of losses team  $i$  has



# Example

$$\begin{bmatrix} 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$



# Generalized Tournament Matrices

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- ▶ Tournament matrix where the values of the entries are 0 and 1 inclusive
- ▶ Entries are the probabilities that one team will defeat another

# Regular Tournament Matrices

A tournament matrix  $A$  of size  $n$  with score vector  $R$  is a regular tournament matrix if

- ▶  $n$  is odd
- ▶ Every entry of  $R$  is  $(n - 1)/2$

Example:

$$\begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

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- ▶ Normal ( $AA^* = A^*A$ )
- ▶ Unitarily Diagonalizable ( $UAU^* =$  diagonal matrix)
- ▶ Spectral radius  $\rho = \rho(A) = (n - 1)/2$
- ▶  $Aj_n = \rho j_n$
- ▶ Tournament matrices of size  $n$  where  $n$  is odd with the largest spectral radius are regular



## Near-Regular Tournament Matrices

A tournament matrix  $A$  of size  $n$  with score vector  $R$  is a near-regular tournament matrix if

- ▶  $n$  is even
- ▶ Half the entries of  $R$  are  $(n - 2)/2$  and the other half are  $n/2$

Example: 
$$\begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

# Construction

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## Theorem

Let  $A$  be any  $n \times n$  tournament matrix. Then,

$$M_A = \begin{bmatrix} A & A^T \\ A^T + I_n & A \end{bmatrix}$$

is a  $2n \times 2n$  near-regular tournament matrix.

## Proof.

Since  $A + A^T = J_n - I_n$ , the first  $n$  rows of  $M_A$  have row sum  $n - 1$  and the last  $n$  rows of  $M_A$  have row sum  $n$ . So the score vector of  $M_A$  is

$$M_{AJ_{2n}} = \begin{bmatrix} (n-1)j_n \\ nj_n \end{bmatrix}.$$

Therefore, by definition,  $M_A$  is a near-regular tournament matrix. □

# Brualdi-Li Matrix

Near-regular tournament matrix of size  $2m$  defined as

$$\mathcal{B}_{2m} = \begin{bmatrix} L_m & L_m^T \\ L_m^T + I_m & L_m \end{bmatrix}$$

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Properties:

- ▶  $\rho(\mathcal{B}_{2m}) \geq \rho(A)$  for every  $2m \times 2m$  tournament matrix  $A$
- ▶ If  $\rho(\mathcal{B}_{2m}) = \rho(A)$ ,  $PAP^T = \mathcal{B}_{2m}$  where  $P$  is some permutation matrix
- ▶ Diagonalizable, though not unitarily
- ▶ First  $m$  entries of score vector are  $n - 1$ . Last  $m$  entries are  $m$ .

# Perron-Frobenius Theorem

## Theorem

*Let  $M$  be a nonnegative, irreducible matrix. Then the spectral radius of  $M$ ,  $\rho(M)$ , is a unique, positive eigenvalue for  $M$ , and there is an entrywise positive eigenvector  $v$ . Such a vector  $v$  is called the Perron vector for  $\rho$ .*

# Kendall-Wei Ranking

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$$\sum_{j=1}^n [A]_{ij} s_j = \sum_{j=1}^n ([A]_{ij} \sum_{k=1}^n [A]_{jk}) = \sum_{k=1}^n \sum_{j=1}^n [A]_{ij} [A]_{jk}$$

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- ▶ This is the sum of all entries in the  $i^{\text{th}}$  row of  $A^2 \rightarrow A^2 j_n$  is the vector whose  $i^{\text{th}}$  entry is the sum of the scores of all teams defeated by team  $i$

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- ▶ Continue process up to  $A^k j_n$  where  $k$  is an arbitrary positive integer pause

- ▶ 
$$\lim_{k \rightarrow \infty} \frac{A^k j_n}{\|A^k j_n\|} = \text{Perron vector } v \text{ (Power Method)}$$

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- ▶ Strength determined by right Perron vector  $v$  ( $Av = \rho v$ )
- ▶ Weakness determined by left Perron vector  $w$  ( $w^T A = \rho w^T$ )

$$\text{▶ } w = \lim_{k \rightarrow \infty} \frac{j_n^T A^k}{\|j_n^T A^k\|}$$

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- ▶  $w = \lim_{k \rightarrow \infty} \frac{j_n^T A^k}{\|j_n^T A^k\|}$

- ▶ Team  $i$  is stronger than team  $j$  if  $v_i/w_i > v_j/w_j$ .



## Brualdi-Li Matrix and Rankings

Let  $\mathcal{B}_{2m}$  be the Brualdi-Li matrix of size  $2m$  with right Perron vector  $v$  and left Perron vector  $w$

- ▶ Kendall-Wei Ranking:

$$v_{2m} < v_{2m-1} < v_{2m-2} < \dots < v_{m+1} < v_1 < v_2 < \dots < v_m$$

- ▶ Ramanujacharyula Ranking:

$$\frac{v_m}{w_m} < \frac{v_1}{w_1} < \frac{v_{m-1}}{w_{m-1}} < \frac{v_2}{w_2} < \frac{v_{m-2}}{w_{m-2}} < \dots < \frac{v_{m/2}}{w_{m/2}} < 1,$$

$$1 < \frac{v_{2m-m/2+1}}{w_{2m-m/2+1}} < \dots < \frac{v_{m+3}}{w_{m+3}} < \frac{v_{2m-1}}{w_{2m-1}} < \frac{v_{2m}}{w_{2m}} < \frac{v_{m+1}}{w_{m+1}}$$

where  $m/2$  is rounded up if  $m$  is odd

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- ▶ Both ranking schemes of  $\mathcal{B}_{2m}$  agree with ranking via score vector

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- ▶ Left Perron vector and right Perron vector are transposes of each other

# Probabilities

Let  $v$  be the Perron vector of a tournament matrix  $A$

- ▶ Probability team  $i$  beats team  $j$  is  $\pi_{ij} = \frac{v_j}{v_i + v_j}$
- ▶ Generalized tournament matrix  $G$ :  $[G]_{ij} = \pi_{ij}$

# Big Example

Consider  $B_{12}$ :

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \end{bmatrix}$$

The right Perron vector is

$$v = \lim_{k \rightarrow \infty} \frac{\mathcal{B}_{12}^k \mathbf{1}_{12}}{\|\mathcal{B}_{12}^k \mathbf{1}_{12}\|} =$$

$$[.282 \quad .279 \quad .275 \quad .269 \quad .261 \quad .250 \quad .296 \quad .298 \quad .302 \quad .307 \quad .313 \quad .323]$$

and the left Perron vector is

$$w = \lim_{k \rightarrow \infty} \frac{\mathbf{1}_{12}^T \mathcal{B}_{12}^k}{\|\mathbf{1}_{12}^T \mathcal{B}_{12}^k\|} =$$

$$[.323 \quad .313 \quad .307 \quad .302 \quad .298 \quad .296 \quad .250 \quad .261 \quad .269 \quad .275 \quad .279 \quad .282]$$

with decimals rounded to three significant figures.

## Strength to weakness ratios

$$\frac{v_6}{w_6} = .845 < \frac{v_1}{w_1} = .873 < \frac{v_5}{w_5} = .876 < \frac{v_2}{w_2} = .890 < \frac{v_4}{w_4} = .891 < \frac{v_3}{w_3} = .896 < 1$$

$$1 < \frac{v_{10}}{w_{10}} = 1.116 < \frac{v_9}{w_9} = 1.122 < \frac{v_{11}}{w_{11}} = 1.123 < \frac{v_8}{w_8} = 1.142 < \frac{v_{12}}{w_{12}} = 1.145 < \frac{v_7}{w_7} = 1.184$$



Ranking according to Kendall-Wei:

12, 11, 10, 9, 8, 7, 1, 2, 3, 4, 5, 6

Ranking according to Ramanucharyula:

7, 12, 8, 11, 9, 10, 3, 4, 2, 5, 1, 6.

## Generalized Tournament Matrix:

0	.503	.506	.512	.519	.530	.488	.486	.483	.479	.474	.466
.497	0	.504	.509	.517	.527	.485	.484	.480	.476	.471	.463
.494	.496	0	.506	.513	.524	.482	.480	.477	.473	.468	.460
.488	.491	.494	0	.507	.518	.476	.474	.471	.467	.462	.454
.481	.483	.487	.493	0	.511	.469	.467	.464	.460	.455	.447
.470	.473	.476	.482	.489	0	.458	.456	.453	.449	.444	.436
.512	.515	.518	.524	.531	.542	0	.498	.495	.491	.486	.478
.514	.516	.520	.526	.533	.544	.502	0	.497	.493	.488	.480
.517	.520	.523	.529	.536	.547	.505	.503	0	.496	.491	.483
.521	.524	.527	.533	.540	.551	.509	.507	.504	0	.495	.487
.526	.529	.532	.538	.545	.556	.514	.512	.509	.505	0	.492
.534	.537	.540	.546	.553	.564	.522	.520	.517	.513	.508	0